

Transactions across Heterogeneous NoSQL Key-Value Data Stores

Akon Dey

The University of Sydney

Motivation

- NoSQL Data Stores
 - Gaining popularity, widely used, sometimes misunderstood
 - Scalable, Distributed, High-performance, Fault-tolerant
- Limited Support for ACID Transactions
 - Single Item or Entity Group transactions
- This works well for some applications
 - When primary-key access and limited scan is sufficient
- But a large number of applications need transactions
 - Megastore, G-Store, Percolator, Spanner, etc. address this
- ACID transactions across multiple items are desired

Also

- Let us also consider other application data
 - State stores, smart device configuration, others
 - This will become more relevant in the very near future
- Why is this even relevant?
 - IPv6 many more devices are uniquely network addressable
 - Networks are faster 4G, Infiniband, fiber to home, . . .
 - Connected devices TVs, Home appliances, etc.
 - Devices are becoming more compute capable
 - ADSL routers, Network printers, others
 - Bonjour (DNS Service Discovery)
- You want to make consistent updates to multiple entities?

Problems

- Data is in Heterogeneous data stores
 - No apriori knowledge of these
- API is different across stores
 - Functionality may be slightly different
 - If-Modified-Since vs. If-None-Match
- There is no cross data store transaction infrastructure
- But the data store supports ...
 - Consistent writes
 - Some form of test-and-set operations
 - Ability to add meta-data to the data record for state

Additionally access privileges to the data store may differ

Observation

- Data stores have single item transactions
 - Atomic operations on single items
 - Conditional update
- Enables moving one record from one consistent state to another
 - Update iff the state of the item is what was originally read
- Idea
 - Use these properties to coordinate transaction across multiple items

Approach

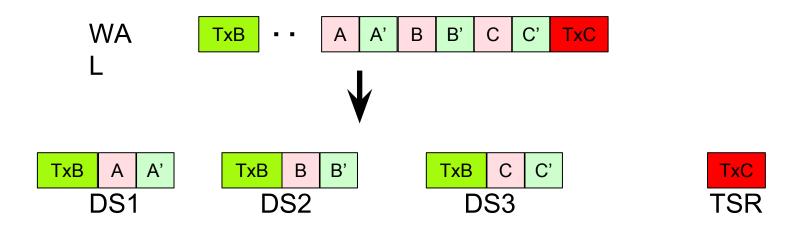
- Break the transaction coordination into two parts
 - Move the coordinator to the client (library) Transaction class
 - Move state information to the data store (meta-data to record)
- Add meta-data to data records
- Global Transaction Status Record (TSR)
 - Globally readable stored in a HA data store
 - Existence indicates transaction is committed

Concurrency Control

- Effectively MVCC
- Similar to Percolator (OSDI'10) ... but
 - No central infrastructure
 - Support heterogeneous data stores
- Use the test-and-set features to implement
- Update records in order of hash(primary key)
 - Global order conflicting transactions has one winner
 - Rest will rollback
- Deadlock avoidance
- Rollback
- Lazy recovery

Logging & Recovery

- Think of it as a deconstructed WAL
 - UNDO + REDO records → Data item
 - Commit record → globally visible Transaction Status Record (TSR)
- Once the TSR exists the transaction is committed
 - If client dies lazy recovery by another client will recover the transaction



Atomic Commit

- We do it in two phases
- Prepare stage: ORDERED test-and-set (ETag or TS)
 - meta-data + data + previous version of data
 - meta-data = TxID, PREPARE, Commit TS, Lease TS
 - If Lease TS expires during prepare stage: rollback (timeout)
 - Write the TSR as COMMITTED
 - Rollforward after this point
- Commit stage: commit all records in parallel
 - Delete the TSR (asynchronously)

Challenges

Time

- Our approach is compatible with TrueTime (time consistency windows)
- TrueTime depends on NTP with atomic and GPS clock at Google scale in a controlled environment
- Reduce network traffic
 - Both the size of messages and the number of messages
- Evaluation
 - Performance
 - Validity (ACID characteristics)

Summary

- Distributed Transactions over autonomous NoSQL data stores
- Prototype Java library works with
 - Google Cloud Storage
 - Window Azure Storage
- Evaluation
 - YSCB with support for Transactions
- Challenges
 - Time
 - Evaluation
- Question and Suggestions
 - akon.dey@sydney.edu.au