OLTP ON NON-VOLATILE MEMORY: WRITE-BEHIND LOGGING

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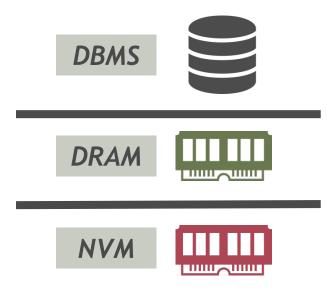
NON-VOLATILE MEMORY (NVM)



Like DRAM, low latency loads and stores

Like SSD, persistent writes and high capacity

CURRENT DATABASE SYSTEMS



Perform write-ahead logging to avoid random writes to durable storage

But, NVM supports fast random writes

TALK OVERVIEW

A new logging algorithm that is better suited for non-volatile memory

Inspired by research done at Wisconsin and Berkeley in the 1990s





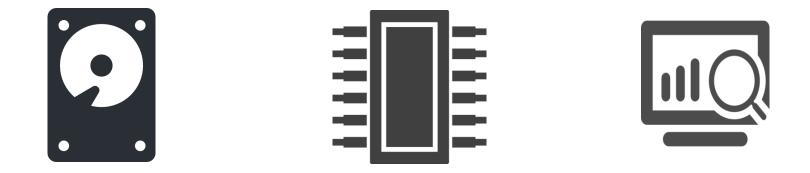
TALK OVERVIEW

A new logging algorithm that is better suited for non-volatile memory

And research done at Puerto Rican DOT(!) and IBM Almaden in the 1970s







WRITE-AHEAD LOGGING

WRITE-BEHIND LOGGING

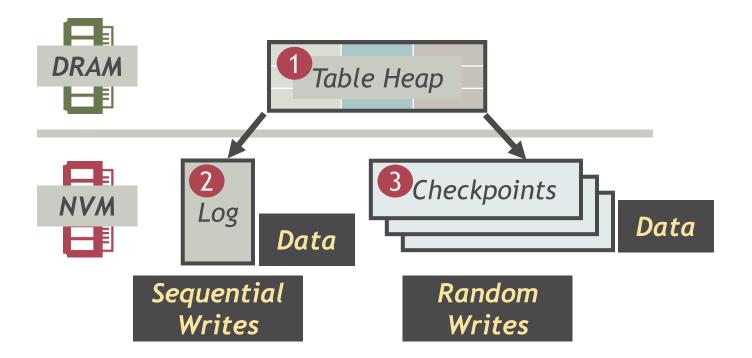
EVALUATION

MULTI-VERSIONED DBMS

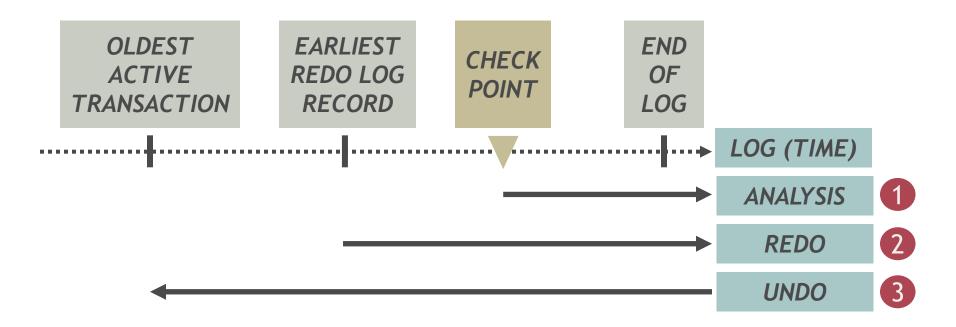
TUPLE ID	BEGIN TIMESTAMP	END TIMESTAMP	PREVIOUS VERSION	TUPLE DATA
1	10	20	-	V-1
2	20	∞	1	V-2



WRITE-AHEAD LOGGING



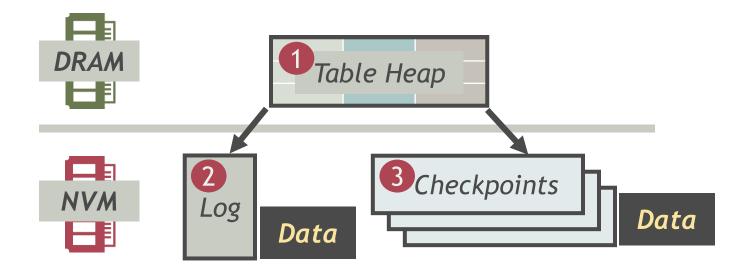
RECOVERY



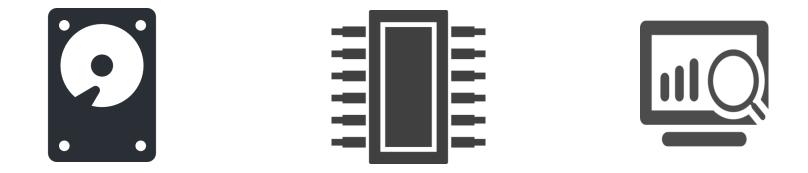
PROBLEM #1: SLOW RECOVERY

- Replaying the redo log takes time
 - Recovery time depends linearly on the number of log records
 - Which in turn depends on frequency of checkpointing
- Slow recovery hurts application availability
 - Even with replication, bringing up new replicas takes time
 - This increases the vulnerability of the system

PROBLEM #2: DATA DUPLICATION



How can we leverage NVM to support faster recovery and reduce data duplication?



WRITE-AHEAD LOGGING

WRITE-BEHIND LOGGING

EVALUATION

WRITE-BEHIND LOGGING

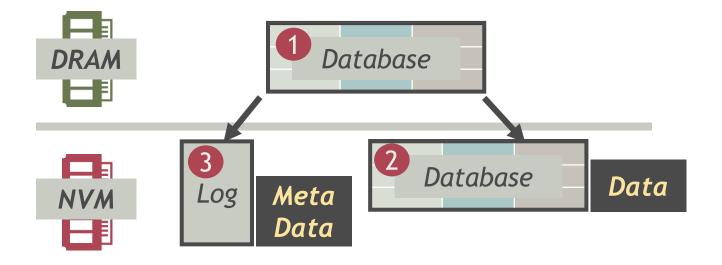
- Write-ahead log serves two purposes
 - Transform random database writes into sequential log writes
 - Support transaction rollback
- Designed for hard disks that can only slow random writes
 - But, NVM supports fast random writes

Why not directly write changes out to the multi-versioned database at runtime?

WRITE-BEHIND LOGGING

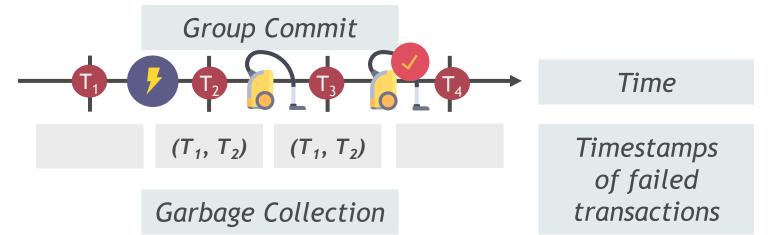
- No physical redo
 - Directly writes changes to the database at runtime
 - Due to multi-versioning, it does not overwrite data
- Logical undo
 - Does not duplicate tuple data in the log and checkpoints
 - Instead, it only records transaction metadata
 - Sufficient to support transaction rollback

WRITE-BEHIND LOGGING



METADATA FOR LOGICAL UNDO

- Record timestamps of failed transactions in log
 - Ignore versions changed during those timestamps (logical undo)



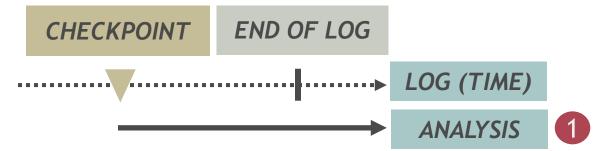
SOLUTION #1: NO DATA DUPLICATION



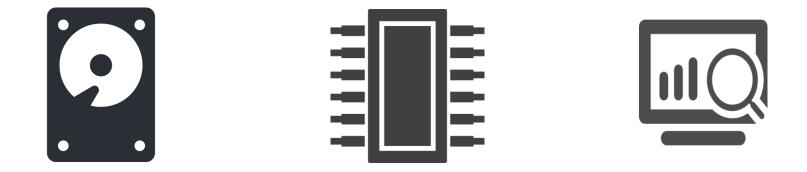
Write-behind logging avoids data duplication by only recording transaction metadata in log

SOLUTION #2: INSTANT RECOVERY

- Single analysis phase to determine failed transactions
 - No physical redo: All changes written to database at runtime
 - Logical undo: Skip reading effects of uncommitted transactions



Write-behind logging enables instant recovery by eliminating redo and doing logical undo



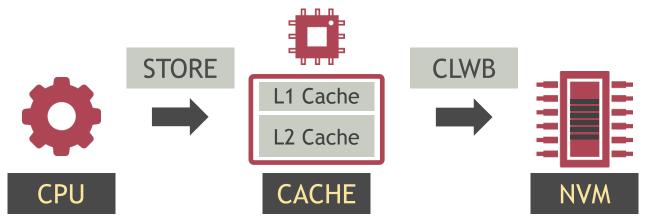
WRITE-AHEAD LOGGING

WRITE-BEHIND LOGGING

EVALUATION

NVM HARDWARE EMULATOR

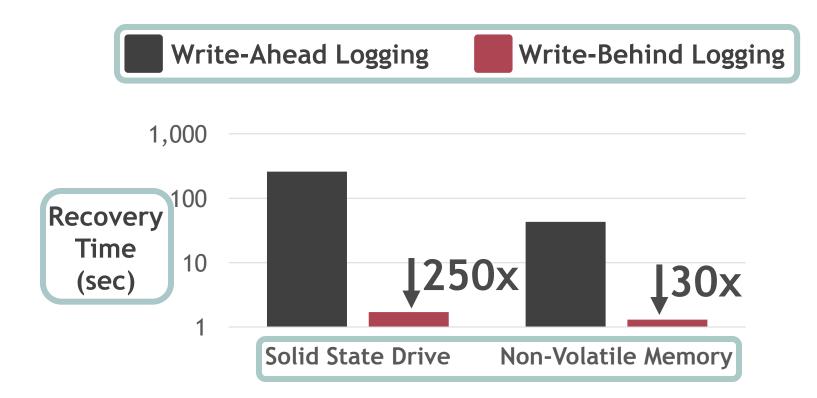
- Special CPU microcode to add stalls on cache misses
 - Tune DRAM latency to emulate different NVM technologies
- New assembly instructions for managing NVM
 - Cache-line write-back (CLWB) instruction



EVALUATION

- Compare logging protocols in Peloton
 - Write-Ahead logging
 - Write-Behind logging
- TPC-C benchmark
- Storage devices
 - Solid-state drive
 - Emulated non-volatile memory

RECOVERY TIME



THROUGHPUT



TAKEAWAYS

- Write-behind logging
 - Enables instant recovery from failures
 - Illustrates importance of rethinking algorithms for NVM
- NVM upends key assumptions about storage
 - This impacts all the layers of a DBMS
 - It's time for complete system rewrite



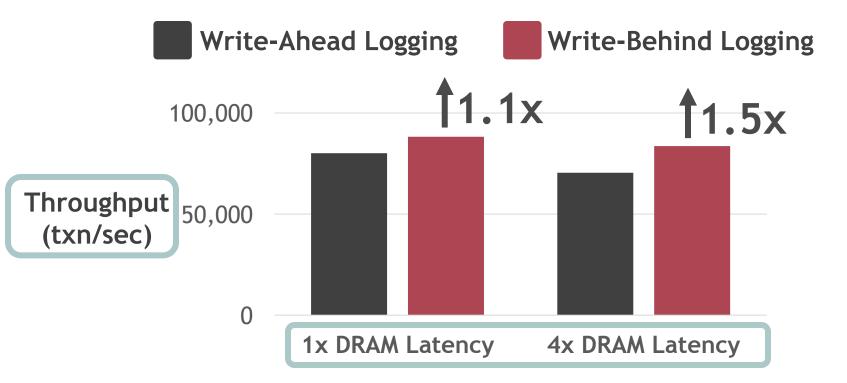
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IMPACT OF NVM LATENCY



REPLICATION

